

A Markov chain for lattice polytopes

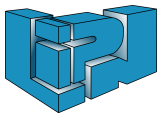
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Outlines

1. Context and motivations
2. Markov Chains
3. Random sampler of (d, k) -polytopes
4. Mixing time
5. Conclusion

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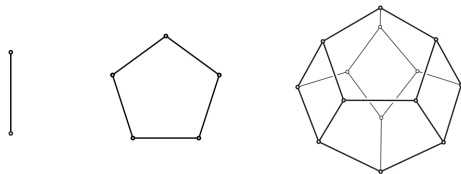
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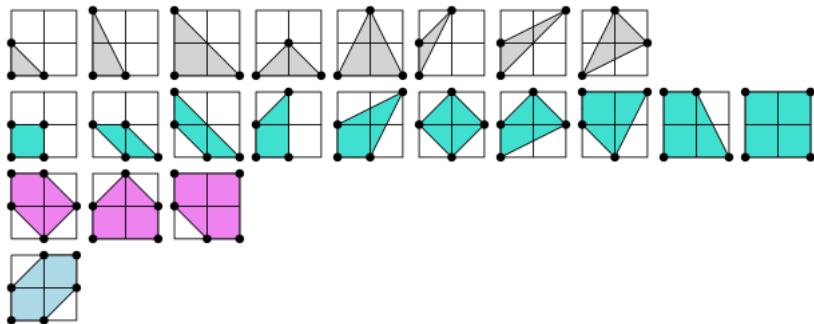
Lattice (d, k) -polytope

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Example $d = 2, k = 2$



Not exhaustive list !!

Motivations

Lattice (d, k) -polytopes :

- ▶ appear in both theoretical and applied mathematical fields.
- ▶ have been studied on their properties as a function of d and k .
- ▶ have a central place in *linear optimization*.

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However :

- ▶ practically few things are known in high dimension : enumeration ? average number of vertices ? ...
- ▶ exhaustive enumeration is prohibited due to combinatoric explosion.

Random sampling

Why?

- ▶ to investigate the average properties of large-sized objects.
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How ?

- ▶ Ad-hoc methods.
- ▶ Combinatoric approach : recursive methods, Boltzmann samplers ...
- ▶ Probabilistic approach : Markov chains.

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What has been done ?

Random samplers when $d = 2$.

- ▶ Random sampler of convex polygons in a disc [Devilleers, Duchon, Thomasse '14]
- ▶ Boltzmann sampler of convex polyominoes [Bodini, Duchon, Jacquot, Mutafchiev '13]

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Our contribution

Uniform random sampler in general dimension using Markov chain.

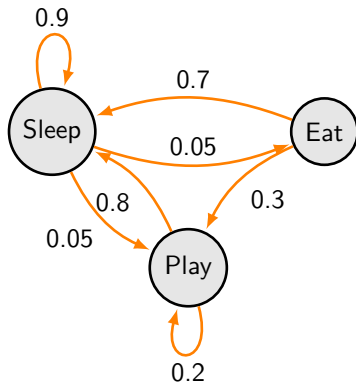
Reminders on Markov Chains

What is Doudou doing?



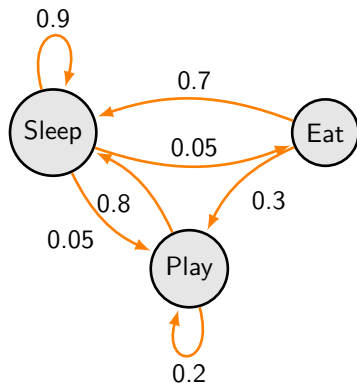
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Reminders on Markov Chains

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Example

The space of states

$$\Omega = \{Sleep, Eat, Play\}$$

(P)

| | Sleep | Eat | Play |
|-------|-------|------|------|
| Sleep | 0.9 | 0.05 | 0.05 |
| Eat | 0.7 | 0 | 0.3 |
| Play | 0.8 | 0 | 0.2 |

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Example

A stationnary distribution for Doudou is :

$$\pi = [0.884 \quad 0.0442 \quad 0.0718].$$

Doudou spends 88.4% of his time sleeping !!

- ▶ Observing Doudou when it reaches its stationnary distribution means we pick at random a state following the distribution π .

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 - ▶ **Note that an for an irreducible chain, every state has the same period.**
 - Symmetry** For any two states x and $y \in \Omega$, $P(x, y) = P(y, x)$.

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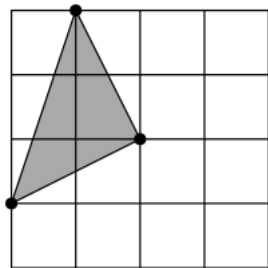
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We want to define the simplest set of transition rules, in order to later facilitate proofs on the necessary length of the walk.

A Markov chain for (d, k) -polytopes

$\Omega :=$ set of lattice (d, k) -polytopes.

Let $P \in \Omega$ be a (d, k) -polytope with vertex set \mathcal{V} .

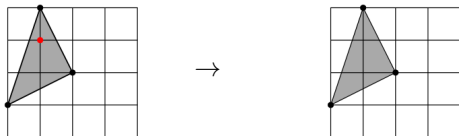


Transition rules

- ▶ sample at random a point $x \in [0, k]^d$.

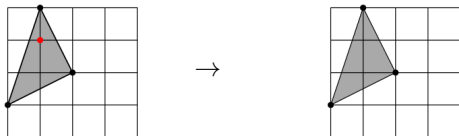
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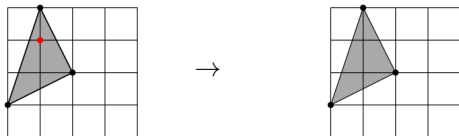


- ▶ If x is a vertex of P , then
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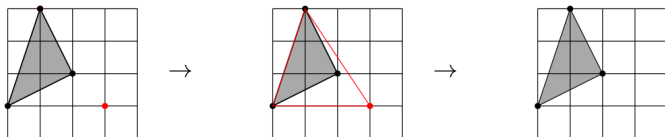


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Corollary

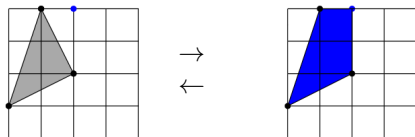
The resulting random sampler is a quasi-uniform random sampler for d -dimensional (d, k) -polytopes.

Resulting random sampler

Sketch of proof

Symmetry Transition rules ensure that for any P and $Q \in \Omega$,

$$\mathbb{P}(P, Q) = \mathbb{P}(Q, P) = \begin{cases} \frac{1}{(k+1)^d} \\ 0 \end{cases}$$

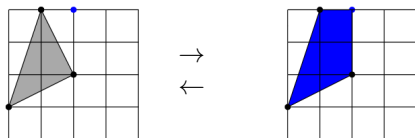


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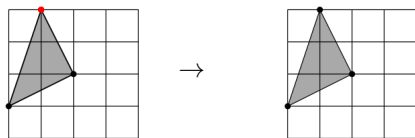
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 3. Given a simplex $S \in \Omega$, we can always perform a succession of insertion and deletion to reach any simplex S' .

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An intuitive lower bound

$$\delta \geq \underbrace{n_{\max}}_{\text{largest number of vertices}} - \underbrace{(d+1)}_{\text{number of vertices of a simplex}} .$$

Lower bounds on mixing time

| dimension | n_{\max} | lower bound on t_{mix} |
|------------|--|--|
| $d = 2$ | $12 \left(\frac{k}{2\pi}\right)^{2/3} + O(k^{1/3} \log k)$ (1) | $t_{\text{mix}} \geq ck^{2/3}$ |
| $d \geq 2$ | $c_1(d)r^d \frac{d-1}{d+1}$ (2) | $t_{\text{mix}} \geq ck^{d \frac{d-1}{d+1}}$ |

(1) Largest number of vertices of a polygon in $[0, k]^2$ [Acketa, Zunic '95, Deza, Manoussakis, Onn '18].

(2) Largest number of faces of each dimension of a lattice polytope contained in a d -dimensional disc of radius r and centered in 0 [Barany, Larman '98].

Experimental results in dimension $d = 2$

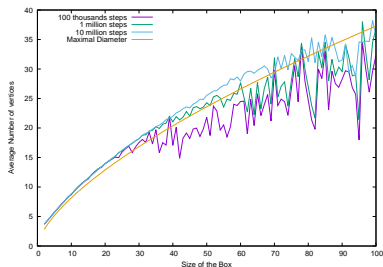


FIGURE – Number of vertices

$$\mathbb{E}[n] \geq 6 \left(\frac{k}{2\pi} \right)^{2/3} \sim \frac{n_{\max}}{2}$$

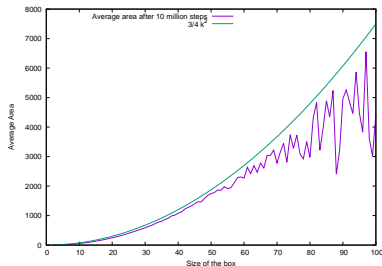


FIGURE – Area

$$\mathbb{E}[a] \leq \frac{3}{4} k^2$$

Conclusion

- ▶ Uniform distribution over the (d, k) -polytopes.
- ▶ Built a random sampler resulting from a Markov chain.
- ▶ Obtain better bounds on mixing time using spectral gap analysis.
- ▶ Find out different transition rules that makes it easier to sample.

Thank you for your attention !